

EXPANSIONS OF THE ORDERED ADDITIVE GROUP OF REAL NUMBERS BY TWO DISCRETE SUBGROUPS

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ABSTRACT. The theory of $(\mathbb{R}, <, +, \mathbb{Z}, \mathbb{Z}a)$ is decidable if a is quadratic. If a is the golden ratio, $(\mathbb{R}, <, +, \mathbb{Z}, \mathbb{Z}a)$ defines multiplication by a . The results are established by using the Ostrowski numeration system based on the continued fraction expansion of a to define the above structures in monadic second order logic of one successor. The converse that $(\mathbb{R}, <, +, \mathbb{Z}, \mathbb{Z}a)$ defines monadic second order logic of one successor, will also be established.

1. INTRODUCTION

Let $a \in \mathbb{R}$. We consider the following structure $\mathcal{R}_a := (\mathbb{R}, <, +, \mathbb{Z}, \mathbb{Z}a)$. Although it is well known that $(\mathbb{R}, <, +, \mathbb{Z})$ has a decidable theory and other desirable model theoretic properties (arguably due to Skolem [13]¹ and later rediscovered independently by Weispfenning [16] and Miller [10]), the question whether the theory of \mathcal{R}_a is decidable even for some irrational number a has been open for a long time. The interest in these structures arises among other things from the observation that the structure \mathcal{R}_a codes many of the Diophantine properties of a . This observation will play a key role throughout this paper. The following is the main result of this paper.

Theorem A. If a is quadratic, then the theory of \mathcal{R}_a is decidable.

A real number is called quadratic if it is the solution to a quadratic equation with rational coefficients. Theorem A provides the first example of an irrational number a such that the theory of \mathcal{R}_a is decidable. Its proof depends crucially on the periodicity of the continued fraction expansion of a . When a is non-quadratic, the conclusion of Theorem A can fail. It will be shown that whenever the continued fraction expansion of a is non-computable, then the theory of \mathcal{R}_a is undecidable. It is also worth noting that while the theory of \mathcal{R}_a can be decidable, its expansion $(\mathbb{R}, <, +, \mathbb{Z}, \mathbb{Z}a, \mathbb{Z}b)$ defines multiplication on \mathbb{R} and hence its theory is undecidable as long as $1, a, b \in \mathbb{R}$ are linearly independent over \mathbb{Q} , by Hieronymi and Tychonievich [8, Theorem C].

Now consider the structure $\mathcal{S}_a := (\mathbb{R}, <, +, \mathbb{Z}, \lambda_a)$, where $\lambda_a : \mathbb{R} \rightarrow \mathbb{R}$ maps x to ax . Note that \mathcal{S}_a is an expansion of \mathcal{R}_a , since $\lambda_a(\mathbb{Z}) = a\mathbb{Z}$. There are more results known about these structures than about \mathcal{R}_a . If a is not a quadratic real number, \mathcal{S}_a defines multiplication on \mathbb{R} and hence its theory is undecidable by [8, Theorem

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¹Skolem essentially showed elimination of quantifiers ranging over elements of \mathbb{Z} . Full quantifier elimination follows easily as pointed out by Smorzyński [14, Exercise III.4.15].

B]. However until now there was no known example of an irrational number a such that the theory of \mathcal{S}_a is decidable. The following Theorem gives the first example of such a real number.

Theorem B. Let $\varphi := \frac{1+\sqrt{5}}{2}$ be the golden ratio. Then \mathcal{R}_φ defines λ_φ and hence the theory of \mathcal{S}_φ is decidable.

Definable here and throughout the paper will always mean definable without parameters. In order to establish Theorem A, we will show that for a quadratic, \mathcal{R}_a is definable in monadic second order logic of one successor. To make this statement precise, consider the two-sorted structure $\mathcal{B} := (\mathbb{N}, \mathcal{P}(\mathbb{N}), s_{\mathbb{N}}, \in)$, where $s_{\mathbb{N}}$ is the successor function on \mathbb{N} and \in is the relation on $\mathbb{N} \times \mathcal{P}(\mathbb{N})$ such that $\in(t, X)$ iff $t \in X$. The structure \mathcal{B} was studied by Büchi in his seminal paper [4]. Using the theory of automata Büchi proved that the theory of \mathcal{B} is decidable and established what would today be called a quantifier elimination result. Theorem A will follow immediately from the decidability of the theory of \mathcal{B} and the following result.

Theorem C. Let $a \in \mathbb{R}$ be quadratic. Then \mathcal{B} defines an isomorphic copy of \mathcal{R}_a .

A structure that is isomorphic to a definable structure in \mathcal{B} is sometimes called Büchi presentable. While Theorem C shows that \mathcal{R}_a is at most as complicated as \mathcal{B} for quadratic a , we will also establish the converse.

Theorem D. Let $a \in \mathbb{R}$ be irrational. Then \mathcal{R}_a defines an isomorphic copy of \mathcal{B} .

One can show that $(\mathbb{R}, <, +, \mathbb{Z})$ does not define an isomorphic copy of \mathcal{B} and is significantly less complicated than \mathcal{B} . Hence Theorem D shows that while the theory of \mathcal{R}_a can be decidable, \mathcal{R}_a is not as well-behaved as $(\mathbb{R}, <, +, \mathbb{Z})$.

To prove Theorems C and D, we will rely on results from the theory of Diophantine approximation. The key tool to construct the isomorphic copies in Theorem C and D will be the Ostrowski representations of both natural numbers and real numbers due to Ostrowski [11]. These representations originating in the theory of Diophantine approximation are based on the continued fraction expansion of a . The reason why the construction in Theorem C works for quadratic numbers and not for others, is that a real number a has a periodic continued fraction expansion if and only if a is quadratic.

This is not the first time that Büchi's Theorem is used to understand expansions of the ordered real additive group. As mentioned by Boigelot, Rassart and Wolper in [3], Büchi himself must have known that the structure $(\mathbb{R}, <, +, \mathbb{Z})$ is definable in \mathcal{B} and hence that its theory is decidable. Also in [3], Büchi's Theorem is used to show that the theory of the expansion of $(\mathbb{R}, <, +, \mathbb{Z})$ by a ternary predicate $V_r(x, u, k)$ that holds iff u is a positive integer power of r , $k \in \{0, \dots, r-1\}$ and the digit of the base- r representation of x in the position corresponding to u is k , is decidable. In some sense, their use of base- r representations will be replaced in this paper by the Ostrowski representations.

The results of the paper should not only be of theoretical importance. The decidability of the theory $(\mathbb{R}, <, +, \mathbb{Z})$ has been used in verification and model checking, since mixed real-integers constraints appear naturally there. Hence the results of

this paper should be relevant in this area, if only by showing that there are interesting expansions of $(\mathbb{R}, <, +, \mathbb{Z})$ whose theory is decidable.

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Notation. We denote $\{0, 1, 2, \dots\}$ by \mathbb{N} . Throughout this paper **definable** will mean definable without parameters.

2. DIOPHANTINE APPROXIMATIONS

In this section we will recall some definitions and results from the study of Diophantine approximations. For more details and proofs, see Rockett and Szűsz [12].

Definition 2.1. A fraction $p/q \in \mathbb{Q}$ is a **best rational approximation** of a real number a if for every fraction p'/q' with $1 \leq q' \leq q$ and $p/q \neq p'/q'$

$$|q'a - p'| > |qa - p|$$

Note that using $|a - p/q|$ instead of $|qa - p|$ changes the definition. For that reason the approximations in previous definition are sometimes called best rational approximation of the second kind.

Definition 2.2. A **continued fraction expansion** $[a_0; a_1, \dots, a_k, \dots]$ is an expression of the form

$$a_0 + \frac{1}{a_1 + \frac{1}{a_2 + \frac{1}{a_3 + \frac{1}{\ddots}}}}$$

For a real number a , we say $[a_0; a_1, \dots, a_k, \dots]$ is the **continued fraction expansion of a** if $a = [a_0; a_1, \dots, a_k, \dots]$ and $a_0 \in \mathbb{Z}$, $a_i \in \mathbb{N}_{>0}$ for $i > 0$.

It is well known that every real number has a unique continued fraction expansion. For the rest of this section, fix a real number a with continued fraction expansion $[a_0; a_1, \dots, a_k, \dots]$.

Fact 2.3. [12, Chapter III.1 Theorem 1 and 2] The continued fraction expansion of a is periodic iff a is a quadratic irrational.

Definition 2.4. Let $k \geq 1$. We define $p_k/q_k \in \mathbb{Q}$ to be the **k -th convergent of a** , that is

$$\frac{p_k}{q_k} = [a_0; a_1, \dots, a_k].$$

The **k -th difference of a** is defined as $\beta_k := q_k a - p_k$.

We define $\zeta_k \in \mathbb{R}$ to be the **k -th complete quotient of a** , that is

$$\zeta_k = [a_k; a_{k+1}, a_{k+2}, \dots].$$

It is worth pointing out that for $k > 0$, $\zeta_k > 1$, since a_k is positive.

Fact 2.5. [12, Chapter II.2 Theorem 2] The set of best rational approximations of a is precisely the set of all convergents of a .

Fact 2.6. [12, Chapter I.1 p. 2] Let $q_{-1} := 0$ and $p_{-1} := 1$. Then $q_0 = 1$, $p_0 = a_0$ and for $k \geq 0$,

$$\begin{aligned} q_{k+1} &= a_{k+1} \cdot q_k + q_{k-1}, \\ p_{k+1} &= a_{k+1} \cdot p_k + p_{k-1}. \end{aligned}$$

It follows immediately that for $k \geq 0$, $\beta_{k+1} = a_{k+1}\beta_k + \beta_{k-1}$.

Fact 2.7. [12, Chapter I.4 p. 9] Let $k \in \mathbb{N}_{>0}$. Then

$$\beta_{k+1} = -\frac{\beta_k}{\zeta_{k+2}}.$$

Since $\zeta_k > 1$, the absolute value of β_k decreases with k .

Fact 2.8. [12, Chapter II.4 p. 24] Let $N \in \mathbb{N}$. Then N can be written uniquely as

$$N = \sum_{k=0}^n b_{k+1}q_k,$$

where $b_k \in \mathbb{N}$ such that $b_1 < a_1$, $b_k \leq a_k$ and, if $b_k = a_k$, $b_{k-1} = 0$.

The representation in the previous fact is called the **Ostrowski representation** of N based on a . This representation will play a crucial role later. If a is the golden ratio, the Ostrowski representation based on a is better known as the Zeckendorf representation, see Zeckendorf [17]. It is important to note that the Ostrowski representation is obtained by a greedy algorithm, see [12, Chapter II.4 p. 24]. The following fact follows immediately.

Fact 2.9. Let $M, N \in \mathbb{N}$ with $M \neq N$ and let $\sum_k b_k q_k$ and $\sum_k c_k q_k$ be the Ostrowski representation of M and N . Let $n \in \mathbb{N}$ be the maximal such that $b_n \neq c_n$. Then $M < N$ iff $b_n < c_n$.

We will also need a similar representation of a real number.

Fact 2.10. [12, Chapter II.6 Theorem 1]² Let $c \in \mathbb{R}$ be such that $-\frac{1}{\zeta_1} \leq c < 1 - \frac{1}{\zeta_1}$. Then c can be written uniquely in the form

$$c = \sum_{k=0}^{\infty} b_{k+1}\beta_k,$$

where $b_k \in \mathbb{N}$, $0 \leq b_1 < a_1$, $0 \leq b_k \leq a_k$, for $k > 1$, and $b_k = 0$ if $b_{k+1} = a_{k+1}$, and $b_k < a_k$ for infinitely many odd k .

One property that is used in the proof of Fact 2.10 is of particular importance to us.

Fact 2.11. [12, Chapter II.6 p.32f] Let $n \in \mathbb{N}$. Then

$$-\beta_n = a_{n+2}\beta_{n+1} + a_{n+4}\beta_{n+3} + a_{n+6}\beta_{n+5} + \dots$$

Hence if n is even, this equation determines the Ostrowski representation of $-\beta_n$.

²While Fact 2.10 is well known, the statement of [12, Chapter II.6 Theorem 1] is unfortunately slightly different. But by inspection of the proof in [12] and using Fact 2.11 and Fact 2.12 the reader can easily verify that Fact 2.10 indeed follows from the statement in [12].

Fact 2.12. Let $n \in \mathbb{N}$ be odd. Then the Ostrowski representation of $-\beta_n$ is

$$\beta_{n-1} + (a_{n+1} - 1)\beta_n + a_{n+3}\beta_{n+2} + a_{n+5}\beta_{n+4} + a_{n+7}\beta_{n+6} + \dots$$

Proof. Since $\beta_{n+1} = a_{n+1}\beta_n + \beta_{n-1}$, we have

$$\begin{aligned} -\beta_n &= (a_{n+1} - 1)\beta_n + \beta_{n-1} - \beta_{n+1} \\ &= \beta_{n-1} + (a_{n+1} - 1)\beta_n + a_{n+3}\beta_{n+2} + a_{n+5}\beta_{n+4} + a_{n+7}\beta_{n+6} + \dots \end{aligned}$$

□

The following fact allows us to decide whether one real number is smaller than another if we are just given their Ostrowski representations.

Fact 2.13. Suppose $\beta_0 > 0$. Let $x, y \in \mathbb{R}$ with $x \neq y$ and let $\sum_k b_{k+1}\beta_k$ and $\sum_k c_{k+1}\beta_k$ be the Ostrowski representations of x and y . Let $n \in \mathbb{N}$ be minimal such that $b_{n+1} \neq c_{n+1}$. Then $x < y$ iff

- (i) $b_{n+1} > c_{n+1}$ and n is odd,
- (ii) $c_{n+1} > b_{n+1}$ and n is even.

Proof. Suppose $b_{n+1} > c_{n+1}$. Then

$$x - y = d_{n+1}\beta_n + d_{n+2}\beta_{n+1} + \sum_{k=n+2}^{\infty} d_{k+1}\beta_k,$$

for some $d_{n+1} > 0$, $d_{n+2} < a_{n+2}$, $-a_k \leq d_k \leq a_k$ for $k \geq n+3$. It is enough show that if n is odd, then $x - y \leq 0$, and if n is even, then $x - y \geq 0$.

Let n be odd. Since $\beta_0 > 0$, we have $\beta_n < 0$ and $\beta_{n+1} > 0$ by Fact 2.7. Since $d_{n+1} \leq a_{n+2} - 1$, we get by Fact 2.11 that

$$\begin{aligned} x - y &\leq \beta_n + (a_{n+2} - 1)\beta_{n+1} + \sum_{k \geq n+2, k \text{ even}} a_{k+1}\beta_k \\ &= \beta_n - \beta_{n+1} + \sum_{k \geq n+1, k \text{ even}} a_{k+1}\beta_k \\ &= -\beta_{n+1} < 0. \end{aligned}$$

Let n be even. Then $\beta_n > 0$ and $\beta_{n+1} < 0$. Since $d_{n+1} \leq a_{n+2} - 1$, we get by Fact 2.11 that

$$\begin{aligned} x - y &\geq \beta_n + (a_{n+2} - 1)\beta_{n+1} + \sum_{k \geq n+2, k \text{ odd}} a_{k+1}\beta_k \\ &= \beta_n - \beta_{n+1} + \sum_{k \geq n+1, k \text{ odd}} a_{k+1}\beta_k \\ &= -\beta_{n+1} > 0. \end{aligned}$$

□

A similar result holds if $\beta_0 < 0$.

Given two natural numbers in Ostrowski representation, it will be important for us to know how to calculate the Ostrowski representation of their sum. Assume for now that a is quadratic. Since the continued fraction expansion of a is periodic, there is a natural number $c := \max_{k \in \mathbb{N}} a_k$. Let $\Sigma_a = \{0, \dots, c\}$ and denote by

Σ_a^* the set of words of finite length on Σ_a . Let $N \in \mathbb{N}$ be such that $\sum_{k=0}^n b_{k+1}q_k$ is the Ostrowski representation of N . Then we define $\rho_a(N)$ to be the Σ_a -word $b_{n+1} \dots b_1$. For $X \subseteq \mathbb{N}^n$ we define

$$0^* \rho_a(X) := \{(0^{l_1} \rho_a(N_1), \dots, 0^{l_n} \rho_a(N_n)) : (N_1, \dots, N_n) \in X, l_1, \dots, l_n \in \mathbb{N}\}.$$

We will now explain what it means for X to be a -recognizable.

Let $\alpha = (\alpha_1, \dots, \alpha_n) \in (\Sigma_a^*)^n$ and let m be the maximal length of $\alpha_1, \dots, \alpha_n$. Then add to each α_i the necessary number of 0's to get a word α'_i of length m . Then the convolution³ of α is defined as the word $\alpha_1 * \dots * \alpha_n \in (\Sigma_a^n)^*$ whose i -th letter is the element of Σ_a^n consisting of the i -th letters of $\alpha'_1, \dots, \alpha'_n$. Let $X \subseteq \mathbb{N}^n$. We say that X is **a -recognizable** if $\{\alpha_1 * \dots * \alpha_n : (\alpha_1, \dots, \alpha_n) \in 0^* \rho_a(X)\} \subseteq (\Sigma_a^n)^*$ is recognizable by a finite automaton. For a definition of a finite automaton and recognizability we refer the reader to Khoussainov and Nerode [9].

Fact 2.14. [7, Theorem B] Let a be quadratic. Then $\{(x, y, z) \in \mathbb{N}^3 : x + y = z\}$ is a -recognizable.

If a is the golden ratio, the Ostrowski representation is called Zeckendorf representation and in this particular case Fact 2.14 was first shown by Frougny in [5]. In [1] Ahlbach et al. present an elementary algorithm to calculate the Zeckendorf representation of a sum in terms of the Zeckendorf representation of the summands. This algorithm was adjusted to give Fact 2.14 by Hieronymi and Terry [7].

3. DEFINING \mathcal{R}_a IN \mathcal{B}

Let $a \in [1, 2]$ be an quadratic irrational number. Since \mathcal{R}_a and \mathcal{R}_{qa} are interdefinable for non-zero $q \in \mathbb{Q}$, we can assume that $1.5 < a < 2$. Let $[a_0; a_1, \dots, a_n, \dots]$ be the continued fraction expansion of a . Note that since $1.5 < a < 2$, $a_0 = a_1 = 1$. By Fact 2.3, the continued fraction expansion of a is periodic. Hence it is of the form

$$[a_0; a_1, \dots, a_\xi, \overline{a_{\xi+1}, \dots, a_{\xi+\nu}}],$$

where ν is the length of the repeating block and the repeating block starts at $\xi + 1$. Set $\mu := \max_i a_i$. Since $a_0 = 1$, we have that β_0 , as is defined in Definition 2.4, is equal to $a - 1$ and hence positive. It also follows easily that $a = 1 + \frac{1}{\zeta_1}$. Hence the interval $[-\frac{1}{\zeta_1}, 1 - \frac{1}{\zeta_1})$ given in Fact 2.10 is equal to the interval $[1 - a, 2 - a)$. We denote this interval by I . Also note that if $\sum_k b_{k+1}\beta_k$ is an Ostrowski representation of a real number in I , then $b_1 = 0$, since $a_1 = 1$. Hence β_0 is not an Ostrowski representation. The same is true for an Ostrowski representation of a natural number. Hence the Ostrowski representation of 1 is q_1 , and not q_0 .

The goal for this section is to show that an isomorphic copy of \mathcal{R}_a is definable in \mathcal{B} . Remember that \mathcal{B} is the two sorted structure $(\mathbb{N}, \mathcal{P}(\mathbb{N}), s_{\mathbb{N}}, \in)$, where $s_{\mathbb{N}}$ is the successor function on \mathbb{N} and \in is the relation on $\mathbb{N} \times \mathcal{P}(\mathbb{N})$ such that $\in(t, X)$ iff $t \in X$. We recall some easy and well-known definability results for \mathcal{B} . We write Even for the set of all even natural numbers and Odd for the set of all odd natural

³Here we followed the presentation in Villemare [15]. For a general definition of convolution see [9].

numbers. Both sets are definable in \mathcal{B} . For example, Even is the unique element X in $\mathcal{P}(\mathbb{N})$ such that

$$0 \in X \wedge \forall x \in \mathbb{N} (x \in X \leftrightarrow s_{\mathbb{N}}(x) \notin X).$$

Similarly, for $m, n \in \mathbb{N}$, the set

$$\{s \in \mathbb{N} : s = m \pmod{n}\}$$

is definable in \mathcal{B} . Also recall that for $m, n \in \mathbb{N}$, we have $m < n$ iff

$$\exists X \in \mathcal{P}(\mathbb{N}) \ m \in X \wedge n \notin X \wedge \forall t \in \mathbb{N} (t \notin X \rightarrow s_{\mathbb{N}}(t) \notin X).$$

Hence the order on \mathbb{N} is definable in \mathcal{B} . If $W \subseteq \mathcal{P}(\mathbb{N})$ is definable in \mathcal{B} , so is the subset W_{fin} of W containing all finite sets in W . Finally, if a subset $X \subset \mathcal{P}(\mathbb{N})^n$ can be recognized by a finite automaton, then it is definable in \mathcal{B} , see for example [4, Lemma 2].

Defining Ostrowski representations. The first step towards defining \mathcal{R}_a in \mathcal{B} will be constructing definable sets that correspond to the Ostrowski representation of both real numbers and natural numbers. This will give us two bijections between definable sets in \mathcal{B} and I and \mathbb{N} .

Definition 3.1. Define $A \subseteq \mathcal{P}(\mathbb{N})^\mu$ to be the set containing $(X_1, \dots, X_\mu) \in \mathcal{P}(\mathbb{N})^\mu$ such that

- $0 \notin X_i$, for $i \geq a_1 - 1$,
- If $n \in X_i$, then $n \notin X_j$ for $j \neq i$,
- $n \notin X_i$, if $0 < n < \xi$ and $i > a_{n+1}$,
- $n \notin X_i$, if $n \geq \xi$, $n + 1 = \xi + l \pmod{\nu}$, $l \in \{1, \dots, \nu\}$ and $i > a_{\xi+l}$,
- for all $m \in \mathbb{N}$ there exists $n \in \text{Even}$ with $n \geq m$ such that there is $l \in \{1, \dots, \nu\}$ with

$$n + 1 = \xi + l \pmod{\nu} \text{ and } n \notin X_{a_{\xi+l}}.$$

It follows from the statements about definability in \mathcal{B} we made before that A is definable in \mathcal{B} . Let $A_{fin} \subseteq A$ be the subset of A containing all tuples (X_1, \dots, X_μ) for which X_i is finite for $i = 1, \dots, \mu$. Since A is definable in \mathcal{B} , so is A_{fin} .

Definition 3.2. Let $b_{k+1} : A_{fin} \rightarrow \mathbb{N}$ map (X_1, \dots, X_μ) to

$$\begin{cases} i, & \text{if } k \in X_i; \\ 0, & \text{otherwise.} \end{cases}$$

If $X \in A_{fin}$, define $Z(X)$ to be the natural number

$$Z(X) = \sum_{k \in \mathbb{N}} b_{k+1}(X) q_k.$$

Note that by uniqueness of Ostrowski representations (see Fact 2.8), the map $Z : A_{fin} \rightarrow \mathbb{N}$ is bijective. Hence Z has an inverse which we denote by Z^{-1} . Also note that the relations $b_{k+1}(X) < b_{k+1}(Y)$ and $b_{k+1}(X) = b_{k+1}(Y)$ on $\mathbb{N} \times A \times A$ are definable in \mathcal{B} .

Definition 3.3. Let $X \in A$. Define $O(X)$ to be the real number in I

$$O(X) = \sum_{k \in \mathbb{N}} b_{k+1}(X) \beta_k.$$

By the uniqueness of the Ostrowski representations (see Fact 2.10), the map $O : A \rightarrow I$ is bijective. Hence O has an inverse which we denote by O^{-1} .

Lemma 3.4. Let $X \in A_{fin}$. Then $aZ(X) - O(X) \in \mathbb{N}$.

Proof. Let $X \in A_{fin}$. Then

$$\begin{aligned} aZ(X) - O(X) &= \sum_{k=0}^n b_{k+1}(X)aq_k - \sum_{k=0}^n b_{k+1}(X)\beta_k \\ &= \sum_{k=0}^n b_{k+1}(X)q_ka - \sum_{k=1}^n b_{k+1}(X)(q_ka - p_k) = \sum_{k=1}^n b_{k+1}(X)p_k \in \mathbb{N}. \end{aligned}$$

□

It is now a good point to outline the strategy for defining \mathcal{R}_a in \mathcal{B} . We have already constructed a bijection O between an interval I and the definable set A in \mathcal{B} . Moreover the map $aZ : A_{fin} \rightarrow \mathbb{Z}a$ that maps $X \in A_{fin}$ to $aZ(X)$ is a bijection. In the following we will amalgamate these two bijection to a single bijection between \mathbb{R} and a set C definable in \mathcal{B} . The reason we choose the map aZ and not the map Z to start with, is Lemma 3.4. Vaguely speaking, because $aZ(X) - O(X) \in \mathbb{N}$, we will be able to recover \mathbb{N} from the images $O(A)$ and $aZ(A_{fin})$.

Defining order and addition. After defining A and A_{fin} , we will now discuss how to define order and addition such that the maps O and Z respect order and addition on I and \mathbb{N} .

Definition 3.5. Let $\oplus : A_{fin} \times A_{fin} \rightarrow A_{fin}$ be given by

$$X \oplus Y := Z^{-1}(Z(X) + Z(Y)).$$

Lemma 3.6. The function \oplus is definable in \mathcal{B} .

Proof. It follows immediately from Fact 2.14 that the graph of \oplus can be recognized by a finite automaton. Hence it is definable in \mathcal{B} . □

Definition 3.7. Let $X, Y \in A_{fin}$ and $X \neq Y$. Let $k \in \mathbb{N}$ be the maximal natural number such that $b_{k+1}(X) \neq b_{k+1}(Y)$. We say $X \prec_Z Y$ if $b_{k+1}(X) < b_{k+1}(Y)$.

It follows immediately from the comment after Definition 3.2 that \prec_Z is definable in \mathcal{B} . The following Lemma follows immediately from Fact 2.9.

Lemma 3.8. Let $X, Y \in A_{fin}$. Then $X \prec_Z Y$ iff $Z(X) < Z(Y)$.

Hence Z is an isomorphism between $(A_{fin}, \prec_Z, \oplus)$ and $(\mathbb{N}, <, +)$.

Lemma 3.9. Let $X, Y \in A_{fin}$. Then

$$O(X \oplus Y) = O(X) + O(Y) \pmod{1}.$$

Proof. By Lemma 3.4, there is $N \in \mathbb{N}$ such that

$$O(X) + O(Y) - O(X \oplus Y) = a(Z(X) + Z(Y) - Z(X \oplus Y)) - N.$$

By definition of \oplus , the right hand side of the previous equation is equal to $-N$. □

Definition 3.10. Let $X, Y \in A$ be such that $X \neq Y$. Let $k \in \mathbb{N}$ be the minimal natural number such that $b_{k+1}(X) \neq b_{k+1}(Y)$. We say $X \prec_O Y$ if one of the following conditions hold:

- (i) $b_{k+1}(X) > b_{k+1}(Y)$ and k is odd,
- (ii) $b_{k+1}(Y) > b_{k+1}(X)$ and k is even.

It is easy to see that $\{(X, Y) \in A : X \prec_O Y\}$ is definable in \mathcal{B} . The following Lemma follows immediately from Fact 2.13.

Lemma 3.11. Let $X, Y \in A$. Then $X \prec_O Y$ iff $O(X) < O(Y)$.

Corollary 3.12. Let $X, Y \in A$ be such that $X \prec_O Y$. Then there is $Z \in A_{fin}$ such that $X \prec_O Z \prec_O Y$.

Proof. By Lemma 3.11, $O(X) < O(Y)$. Take $Z_0 \in A$ such that $O(X) < O(Z_0) < O(Y)$. Let $\varepsilon \in \mathbb{R}_{>0}$ be such that $\varepsilon < \min\{O(Z_0) - O(X), O(Y) - O(Z_0)\}$. Take $n \in \mathbb{N}$ such that

$$\max\left\{\sum_{k>n, k \text{ even}} a_{k+1}\beta_k, -\sum_{k\geq n, k \text{ odd}} a_{k+1}\beta_k\right\} < \varepsilon.$$

Let $Z_{0,1}, \dots, Z_{0,\mu} \in \mathcal{P}(\mathbb{N})$ such that $Z_0 = (Z_{0,1}, \dots, Z_{0,\mu})$. Set

$$Z := ((Z_{0,1} \cap [0, n], \dots, Z_{0,\mu} \cap [0, n]).$$

Then

$$\begin{aligned} O(Z) &= \sum_{k=0}^n b_{k+1}(Z)\beta_k = \sum_{k=0}^{\infty} b_{k+1}(Z_0)\beta_k - \sum_{k=n+1}^{\infty} b_{k+1}(Z)\beta_k \\ &= O(Z_0) - \sum_{k=n+1}^{\infty} b_{k+1}(Z_0)\beta_k. \end{aligned}$$

Since $|\sum_{k=n+1}^{\infty} b_{k+1}(Z_0)\beta_k| < \varepsilon$, we have $O(X) < O(Z) < O(Y)$. Again by Lemma 3.11, $X \prec_O Z \prec_O Y$. \square

We now use this density to extend \oplus to A .

Definition 3.13. Define $+_1 : I \times I \rightarrow I$ be the function that maps $(c_1, c_2) \in I^2$ to the unique element $d \in I$ such that $d = c_1 + c_2 \pmod{1}$. Let $\oplus : A \times A \rightarrow A$ be function that maps (X, Y) to $O^{-1}(O(X) +_1 O(Y))$.

The following Lemma follows immediately from the definition of \oplus .

Lemma 3.14. Let $X, Y \in A$. Then

$$O(X \oplus Y) = O(X) + O(Y) \pmod{1}.$$

Lemma 3.15. The map $\oplus : A \times A \rightarrow A$ is definable in \mathcal{B} .

Proof. Consider the following two structures,

$$\mathcal{M}_1 := (A, \prec_O, \oplus|_{A_{fin}}, A_{fin}), \mathcal{M}_2 := (I, <, +_1|_{O(A_{fin})}, O(A_{fin})).$$

By Lemma 3.9 and Lemma 3.11, the map $O : \mathcal{M}_1 \rightarrow \mathcal{M}_2$ is an isomorphism. Let \mathcal{T} be the topology on I whose basic open sets are the intervals (c_1, c_2) , if $c_1, c_2 \in I$ and $c_1 < c_2$, and the sets $[1 - a, c_2) \cup (c_1, 2 - a)$, if $c_1, c_2 \in I$ and $c_1 > c_2$. It is immediate that the topological closure with respect to \mathcal{T} of a set definable in \mathcal{M}_2 is definable in \mathcal{M}_2 as well. By Lemma 3.11 and Lemma 3.12, $O(A_{fin})$ is dense in I with respect to \mathcal{T} . Because of the continuity of $+_1$ with respect to \mathcal{T} , the topological closure in \mathcal{T} of the graph $+_1|_{O(A_{fin})}$ is the graph of $+_1$. Hence $+_1$ is definable in \mathcal{M}_2 . Since O is an isomorphism, \oplus is definable in \mathcal{M}_1 and hence in \mathcal{B} . \square

Thus (A, \oplus) forms a group. The neutral element $\mathbf{0}$ is $(\emptyset, \dots, \emptyset)$. We write

$$\mathbf{1} := (\{1\}, \emptyset, \dots, \emptyset).$$

Note that $O(\mathbf{1}) = \beta_1 = a - 2$. For $X \in A$, we denote the inverse of X with respect to \oplus by $\ominus X$, that means $X \oplus (\ominus X) = \mathbf{0}$. As usual, for $X, Y \in A$ we will write $X \ominus Y$ for $X \oplus (\ominus Y)$.

Modifying O . We have constructed an isomorphism O between (A, \prec_O, \oplus) and $(I, <, +_1)$. In the following this isomorphism will be modified to an isomorphism S whose range is $([0, 1], <, + \bmod 1)$ instead of $(I, <, +_1)$. Here $+ \bmod 1 : [0, 1]^2 \rightarrow [0, 1]$ is the map that takes $(x, y) \in [0, 1]^2$ to the unique $z \in [0, 1]$ such that $x + y = z \bmod 1$.

Definition 3.16. Let $X, Y \in A$, we write $X \prec_1 Y$ if

- $Y \prec_O \mathbf{1} \preceq_O X$,
- $X, Y \prec_O \mathbf{1}$ and $X \prec_O Y$,
- $X, Y \succ_O \mathbf{1}$ and $X \prec_O Y$

Let $\oplus_1 : A \times A \rightarrow A$ be the map that takes $(X, Y) \in A^2$ to $(X \oplus Y) \ominus \mathbf{1}$. Let $S : A \rightarrow [0, 1]$ maps $X \in A$ to

$$\begin{cases} O(X) - O(\mathbf{1}), & \text{if } \mathbf{1} \preceq_O X; \\ O(X) - O(\mathbf{1}) + 1, & \text{otherwise.} \end{cases}$$

Lemma 3.17. The map $S : (A, \prec_1, \oplus_1) \rightarrow ([0, 1], <, + \bmod 1)$ is an isomorphism.

Proof. Since $O(\mathbf{1}) = a - 2$ and $I = [1 - a, 2 - a)$, we directly get that

$$S(\{X \in A : \mathbf{1} \preceq_O X\}) = [0, 4 - 2a) \text{ and } S(\{X \in A : X \prec_O \mathbf{1}\}) = [4 - 2a, 1).$$

We have $4 - 2a < 1$, since $1.5 < a < 2$. It follows immediately that $S(X) < S(Y)$ iff $X \prec_1 Y$. Since O is bijective, it is easy to see that S is bijective. Note that $S(X)$ is the unique $c \in [0, 1]$ with $c = O(X) - a \bmod 1$. Let $X, Y \in A$. Then

$$\begin{aligned} S(X) + S(Y) &= O(X) + O(Y) - 2O(\mathbf{1}) \bmod 1 \\ &= O((X \oplus Y) \ominus \mathbf{1}) - O(\mathbf{1}) \bmod 1 \\ &= S(X \oplus_1 Y) \bmod 1. \end{aligned}$$

□

Hence (A, \oplus_1) is a group and its neutral element is $\mathbf{1}$. As above, for $X \in A$ we will write $\ominus_1 X$ for the inverse element of X in A with respect to \oplus_1 . Thus $\ominus_1 X$ is the unique element in A such that $(\ominus_1 X) \oplus_1 X = \mathbf{1}$.

Lemma 3.18. Let $X, Y \in A$. Then $\ominus_1 X \preceq_1 Y$ iff $S(X) + S(Y) \geq 1$.

Proof. Since $1 - S(X) = -S(X) \bmod 1$ and S is group homomorphism, we have $1 - S(X) = S(\ominus_1 X)$. Hence we have by Lemma 3.17

$$\begin{aligned} S(X) + S(Y) \geq 1 &\text{ iff } S(Y) \geq 1 - S(X) \\ &\text{ iff } S(Y) \geq S(\ominus_1 X) \\ &\text{ iff } Y \succeq_1 \ominus_1 X. \end{aligned}$$

□

Corollary 3.19. Let $X, Y \in A$. Then

$$S(X \oplus_1 Y) = \begin{cases} S(X) + S(Y), & \text{if } \ominus_1 X \preceq_1 Y; \\ S(X) + S(Y) - 1, & \text{otherwise.} \end{cases}$$

Recovering \mathbb{N} . We have now established that we can define order and addition on A and A_{fin} such that O , S and Za become isomorphisms. Vaguely speaking, the next step is to recover \mathbb{N} from $O(A)$ and $Za(A_{fin})$. We will find a set B definable in \mathcal{B} , a definable order \prec_B on B , a definable operation $\oplus_B : B \times B \rightarrow B$ and a map $R : B \rightarrow \mathbb{N}$ such that R is an isomorphism between (B, \prec_B, \oplus_B) and $(\mathbb{N}, <, +)$. It will be crucial later that the isomorphism R arises naturally from O and Za .

Lemma 3.20. Let $X \in A_{fin}$. Then $\mathbb{N} \cap (aZ(X), (Z(X) + 1)a)$ is

$$\begin{cases} \{aZ(X) - O(X) + 1, aZ(X) - O(X) + 2\}, & \text{if } X \prec_O \mathbf{1} \\ \{aZ(X) - O(X) + 1\}, & \text{if } \mathbf{1} \preceq_O X \preceq_O \mathbf{0}; \\ \{aZ(X) - O(X), aZ(X) - O(X) + 1\}, & \text{otherwise.} \end{cases}$$

Proof. Let $X \in A_{fin}$. Since $1.5 < a < 2$, there are at most two natural numbers between $aZ(X)$ and $a(Z(X) + 1)$. Moreover, since $1.5 < a < 2$, $3 - a < a$. We will use this observation repeatedly throughout this proof. By Lemma 3.4, $aZ(X) - O(X) \in \mathbb{N}$ and so are $aZ(X) - O(X) + 1$ and $aZ(X) - O(X) + 2$. Since $1 - a \leq O(X) < 2 - a$, we also have that

$$aZ(X) < aZ(X) + O(X) + 1 < a(Z(X) + 1).$$

We just have to determine which of the other two natural numbers fall into the interval we are considering. First consider the case that $X \prec_O \mathbf{1}$. Since $O(\mathbf{1}) = a - 2$, we have $O(X) < a - 2$ by Lemma 3.11. Since $1 - a \leq O(X) < a - 2$, we get that

$$aZ(X) + O(X) < aZ(X) < aZ(X) + O(X) + 2 < a(Z(X) + 1).$$

Now suppose that $\mathbf{1} \preceq_O X \preceq_O \mathbf{0}$. Since $O(\mathbf{0}) = 0$ and $O(\mathbf{1}) = a - 2$, we have $a - 2 \leq O(X) \leq 0$ by Lemma 3.11. Hence

$$aZ(X) + O(X) \leq aZ(X) \text{ and } a(Z(X) + 1) \leq aZ(X) + O(X) + 2.$$

Finally consider the case that $\mathbf{0} \prec_O X$. Since $O(\mathbf{0}) = 0$, we have $0 < O(X)$ by Lemma 3.11. Since $O(X) < 2 - a$, we have

$$aZ(X) < aZ(X) + O(X) < aZ(X) + O(X) + 1 < a(Z(X) + 1).$$

□

Definition 3.21. Let $B \subseteq A_{fin} \times \{0, 1, 2\}$ be defined as the set of all pairs (X, i) that have one of the following properties:

- (i) $X = \mathbf{0}$ and $i = 0$,
- (ii) $X \prec_O \mathbf{1}$ and $i \in \{1, 2\}$,
- (iii) $\mathbf{1} \preceq_O X \preceq_O \mathbf{0}$ and $i = 1$,
- (iv) $\mathbf{0} \prec_O X$ and $i \in \{0, 1\}$.

Let $R : B \rightarrow \mathbb{N}$ map (X, i) to $aZ(X) - O(X) + i$.

Lemma 3.22. R is a bijection.

Proof. By Lemma 3.20, $\mathbb{N}a \cap \mathbb{N} = \{0\}$ and the fact that $Z : A_{fin} \rightarrow \mathbb{N}$ is a bijection, R maps $B \setminus \{(\mathbf{0}, 0)\}$ bijectively to $\mathbb{N}_{>0}$. Hence R is bijective, since $R((\mathbf{0}, 0)) = 0$. □

Definition 3.23. Let $(X, i), (Y, j) \in B$. We write $(X, i) \prec_B (Y, j)$ if either

- $X = Y$ and $i < j$ or
- $X \prec_Z Y$.

Let $s_B : B \rightarrow B$ map $Z \in B$ to its \prec_B -successor in B . Let $p_B : B \setminus \{(\mathbf{0}, 0)\} \rightarrow B$ map $Z \in B$ to its \prec_B -predecessor in B .

Since \prec_Z well-orders A_{fin} , \prec_B well-orders B . Hence the successor and predecessor function are well-defined. Moreover, by Lemma 3.8, we have that for $Z_1, Z_2 \in B$, $Z_1 \prec_B Z_2$ iff $R(Z_1) < R(Z_2)$. Since R is a bijection, we have for $Z \in B$

$$(3.1) \quad R(s_B(Z)) = R(Z) + 1.$$

We will use the following notation: we write s_B^0 for the identity on B , and for $i \in \mathbb{N}_{>0}$, we write s_B^{-i} for i -th iterate of p_B and s_B^i for the i -th iterate of s_B .

We will now define $\oplus_B : B \times B \rightarrow B$ such that R is an isomorphism from (B, \prec_B, \oplus) to $(\mathbb{N}, <, +)$. For convenience set $\mathbf{e} := \ominus \mathbf{1}$. Since $O(\mathbf{1}) = a - 2$, we have that $O(\mathbf{e}) = 1 - a$ by Lemma 3.14. Hence \mathbf{e} is the left endpoint of I and is \prec_O -minimal in A by Lemma 3.11.

Definition 3.24. For $X, Y \in A$, we define $r(X, Y) \in \{0, 1, 2\}$ to be

$$\begin{cases} 0, & \text{if } X \prec_O \mathbf{e} \ominus Y \text{ and } Y \preceq_O \mathbf{0}; \\ 2 & \text{if } \mathbf{e} \ominus Y \preceq_O X \text{ and } Y \succ_O \mathbf{0}; \\ 1, & \text{otherwise.} \end{cases}$$

Definition 3.25. Let $(X, i), (Y, j) \in B$. We define $\oplus_B : B \times B \rightarrow B$ by

$$(X, i) \oplus_B (Y, j) := s^{i+j+r(X,Y)-2}((X \oplus Y, 1)).$$

Lemma 3.26. Let $X \in A$. Then

$$O(\mathbf{e} \ominus X) = \begin{cases} O(\mathbf{e}) - O(X), & \text{if } X \preceq_O \mathbf{0}; \\ O(\mathbf{e}) - O(X) + 1, & \text{otherwise.} \end{cases}$$

Proof. Suppose $X \preceq_O \mathbf{0}$. Since $1 - a \leq O(X) \leq 0$, we have

$$1 - a \leq O(\mathbf{e}) - O(X) = 1 - a - O(X) \leq 0.$$

Hence $O(\mathbf{e}) - O(X) \in I$ and thus $O(\mathbf{e} \ominus X) = O(\mathbf{e}) - O(X)$ by Lemma 3.14.

Suppose $X \succ_O \mathbf{0}$. Then $0 < O(X) < 2 - a$, and thus $0 < a - 1 < 1 - O(X) < 1$. Hence

$$O(\mathbf{e}) < O(\mathbf{e}) - O(X) + 1 < 1 - a + 1 = 2 - a.$$

Hence $O(\mathbf{e}) - O(X) + 1 \in I$ and therefore $O(\mathbf{e} \ominus X) = O(\mathbf{e}) - O(X) + 1$. \square

Lemma 3.27. Let $X, Y \in A$. Then $O(X) + O(Y) = O(X \oplus Y) + r(X, Y) - 1$.

Proof. We first consider the case that $Y \preceq_O \mathbf{0}$. Suppose that $X \prec_O \mathbf{e} \ominus Y$. Hence $r(X, Y) = 0$. By Lemma 3.26

$$O(X) + O(Y) < O(\mathbf{e} \ominus Y) + O(Y) = O(\mathbf{e}) \leq O(X \oplus Y).$$

Hence $O(X) + O(Y) + 1 \in I$ and thus $O(X \oplus Y) - 1 = O(X) + O(Y)$. Suppose now that $X \succeq_O \mathbf{e} \ominus Y$. Hence $r(X, Y) = 1$. By Lemma 3.26

$$O(X) \geq O(X) + O(Y) \geq O(\mathbf{e} \ominus Y) + O(Y) = O(\mathbf{e}).$$

Hence $O(X) + O(Y) \in I$ and thus $O(X \oplus Y) = O(X) + O(Y)$.

Now consider the case that $Y \succ_O \mathbf{0}$. Suppose that $\mathbf{e} \oplus Y \preceq_O X$. Hence $r(X, Y) = 2$. Again by Lemma 3.26

$$O(X) + O(Y) \geq O(Y) + O(\mathbf{e}) - O(Y) + 1 = O(\mathbf{e}) + 1 = 2 - a > O(X \oplus Y).$$

Thus $O(X) + O(Y) = O(X \oplus Y) + 1$. Finally suppose that $\mathbf{e} \oplus Y \succ_O X$. Then we have $r(X, Y) = 1$. By Lemma 3.26

$$O(X) \leq O(X) + O(Y) < O(Y) + O(\mathbf{e}) + 1 - O(Y) = O(\mathbf{e}) + 1 = 2 - a.$$

Hence $O(X) + O(Y) \in I$ and hence $O(X \oplus Y) = O(X) + O(Y)$. \square

Lemma 3.28. Let $Z_1, Z_2 \in B$. Then $R(Z_1 \oplus_B Z_2) = R(Z_1) + R(Z_2)$.

Proof. Let $(X, i), (Y, j) \in B$. Then by (3.1) and Lemma 3.27

$$\begin{aligned} R((X, i) \oplus_B (Y, j)) &= R(s^{i+j+r(X, Y)-2}(X \oplus Y, 1)) \\ &= R((X \oplus Y), 1) + i + j + r(X, Y) - 2 \\ &= Z(X \oplus Y)a + O(X \oplus Y) + i + j + r(X, Y) - 1 \\ &= aZ(X) + Z(Y)a + O(X) + O(Y) + i + j \\ &= R((X, i)) + R((Y, j)). \end{aligned}$$

\square

Corollary 3.29. The map $R : (B, \prec_B, \oplus_B) \rightarrow (\mathbb{N}, <, +)$ is an isomorphism.

Amalgamating R and S . We have constructed two isomorphisms $R : B \rightarrow \mathbb{N}$ and $S : A \rightarrow [0, 1)$. We define $T : B \times A \rightarrow \mathbb{R}_{\geq 0}$ as the map that takes $(Z, X) \in B \times A$ to $R(Z) + S(X)$. It follows immediately from Lemma 3.17 and Lemma 3.22 that T is bijective. We will now construct two definable subsets A', B' of $B \times A$, a definable relation \prec_C and a definable operation $\oplus_C : (B \times A)^2 \rightarrow B \times A$ such that T is an isomorphism between $(B \times A, \prec_C, \oplus_C, B', A')$ and $(\mathbb{R}_{\geq 0}, <, +, \mathbb{N}, \mathbb{N}a)$.

Definition 3.30. Set $C := B \times A$. Let $A' \subseteq C$ be

$$\begin{aligned} &\{(p_B^2(X, 1), X \oplus \mathbf{1}) : X \in A_{fin}, X \prec_O \mathbf{0}\} \\ &\cup \{(p_B(X, 1), X \oplus \mathbf{1}) : X \in A_{fin}, X \succeq_O \mathbf{0}\}, \end{aligned}$$

and let $B' \subseteq C$ be the set $\{(Z, \mathbf{1}) : Z \in B\}$.

Lemma 3.31. The map $T : (C, B', A') \rightarrow (\mathbb{R}_{\geq 0}, \mathbb{N}, \mathbb{N}a)$ is an isomorphism.

Isomorphism	Domain	Codomain
O	A	I
Z	A_{fin}	\mathbb{N}
S	A	$[0, 1)$
R	B	\mathbb{N}
T	C	$\mathbb{R}_{\geq 0}$

TABLE 1. A list of the isomorphisms and their domain and codomain

Proof. We first show that $T(B') = \mathbb{N}$. Let $(Z, \mathbf{1}) \in B'$. Then

$$T(Z, \mathbf{1}) = R(Z) + S(\mathbf{1}) = R(Z) + 0 = R(Z) \in \mathbb{N}.$$

Since $R : B \rightarrow \mathbb{N}$ is bijective by Lemma 3.22, we have $T(B') = \mathbb{N}$.

We now establish that $T(A') = \mathbb{N}a$. Let $X \in A_{fin}$. Then

$$(3.2) \quad O(X \oplus \mathbf{1}) = O(X) + O(\mathbf{1}) \pmod{1}.$$

Suppose that $X \succeq_O \mathbf{0}$. Since $0 \leq O(X) < 1$ and $S(X \oplus \mathbf{1}) \in [0, 1)$, we get $S(X \oplus \mathbf{1}) = O(X)$ by (3.2). Hence

$$\begin{aligned} T(p_B(X, 1), X \oplus \mathbf{1}) &= R(p_B(X, 1)) + S(X \oplus \mathbf{1}) \\ &= aZ(X) - O(X) + O(X) = aZ(X). \end{aligned}$$

Now suppose $X \prec_O \mathbf{0}$. Since $0 < O(X) + 1 < 1$, we have $S(X \oplus \mathbf{1}) = O(X) + 1$ by (3.2). Hence

$$\begin{aligned} T(p_B^2(X, 1), X \oplus \mathbf{1}) &= R(p_B^2(X, 1)) + S(X \oplus \mathbf{1}) \\ &= aZ(X) - O(X) - 1 + O(X) + 1 = aZ(X). \end{aligned}$$

Since $Z : A_{fin} \rightarrow \mathbb{N}$ is bijective, we have $T(A') = a\mathbb{N}$. □

Definition 3.32. Let $(Z_1, X_1), (Z_2, X_2) \in C$, we define

$$(Z_1, X_1) \oplus_C (Z_2, X_2) := \begin{cases} (s_B(Z_1 \oplus_B Z_2), X_1 \oplus_1 X_2), & \text{if } \ominus_1 X_1 \preceq_1 X_2; \\ (Z_1 \oplus_B Z_2, X_1 \oplus_1 X_2), & \text{otherwise.} \end{cases}$$

We say $(Z_1, X_1) \prec_C (Z_2, X_2)$ if $Z_1 \prec_B Z_2$ or $(Z_1 = Z_2 \text{ and } X_1 \prec_1 X_2)$.

Lemma 3.33. The map $T : (C, \prec_C, \oplus_C, B', A') \rightarrow (\mathbb{R}_{\geq 0}, <, +, \mathbb{N}, \mathbb{N}a)$ is an isomorphism.

Proof. Let $(Z_1, X_1), (Z_2, X_2) \in C$. By the definition of the maps R and S , we have $R(B) = \mathbb{N}$ and $S(A) = [0, 1)$. Thus we see directly that $T(Z_1, X_1) < T(Z_2, X_2)$ holds iff either $R(Z_1) < R(Z_2)$ holds or, $R(Z_1) = R(Z_2)$ and $S(Z_1) < S(Z_2)$ hold. By Corollary 3.29 and Lemma 3.17 we have $T(Z_1, X_1) < T(Z_2, X_2)$ iff $(Z_1, X_1) \prec_C (Z_2, X_2)$.

It is left to show that $T((Z_1, X_1) \oplus_C (Z_2, X_2)) = T(Z_1, X_1) + T(Z_2, X_2)$. First suppose that $\ominus_1 X_1 \preceq_1 X_2$. Then by Corollary 3.19

$$\begin{aligned} T((Z_1, X_1) \oplus_C (Z_2, X_2)) &= T(s_B(Z_1 \oplus_B Z_2), X_1 \oplus_1 X_2) \\ &= R(s_B(Z_1 \oplus_B Z_2)) + S(X_1 \oplus_1 X_2) \\ &= R(Z_1 \oplus_B Z_2) + 1 + S(X_1) + S(X_2) - 1 \\ &= R(Z_1) + R(Z_2) + S(X_1) + S(X_2) \\ &= T(Z_1, X_1) + T(Z_2, X_2). \end{aligned}$$

If $\ominus_1 X_1 \succ_1 X_2$, we have by Corollary 3.19

$$\begin{aligned}
T((Z_1, X_1) \oplus_C (Z_2, X_2)) &= T(Z_1 \oplus_B Z_2, X_1 \oplus_1 X_2) \\
&= R(Z_1 \oplus_B Z_2) + S(X_1 \oplus_1 X_2) \\
&= R(Z_1 \oplus_B Z_2) + S(X_1) + S(X_2) \\
&= T(Z_1, X_1) + T(Z_2, X_2).
\end{aligned}$$

□

It follows easily from the previous Lemma that an isomorphic copy of \mathcal{R}_a is definable in \mathcal{B} . Hence Theorem C holds.

4. DEFINING \mathcal{B} IN \mathcal{R}_a

Let $a \in \mathbb{R} \setminus \mathbb{Q}$. Since \mathcal{R}_a and \mathcal{R}_{qa} are interdefinable for non-zero $q \in \mathbb{Q}$, we can assume that $1.5 < a < 2$. In this section, we will show that an isomorphic copy of \mathcal{B} is definable in \mathcal{R}_a . We do not require a to be quadratic.

Since $1 < a < 2$, we have $a = 1 + \frac{1}{\zeta_1}$ and hence $[1 - a, 2 - a) = [-\frac{1}{\zeta_1}, 1 - \frac{1}{\zeta_1})$. Recall that we denote this interval I . It is obviously definable in \mathcal{R}_a . Moreover, since $1 < a < 2$, $\beta_0 = a - 1 > 0$.

Definition 4.1. Let U be the set of all pairs $(p, qa) \in \mathbb{N} \times \mathbb{N}a$ with

$$\forall p' \in \mathbb{N} \forall q'a \in \mathbb{N}a (a \leq q'a \leq qa \wedge (p, qa) \neq (p', q'a)) \rightarrow |q'a - p'| > |qa - p|$$

Note that U is definable in \mathcal{R}_a . By Fact 2.5 the set $\{q_k a : k > 0\}$ is the projection on the second coordinate of U and hence definable in \mathcal{R}_a . We denote this set by V . Since V is definable, the successor function s_V on V is definable as well. Note for every $q_l a \in V$ we have $s_V(q_l a) = q_{l+1} a$.

Definition 4.2. Let $f : \mathbb{N}a \rightarrow \mathbb{R}$ map na to $na - m$, where m is the unique natural number such that $na - m \in I$.

Obviously, f is well-defined and definable in \mathcal{R}_a .

Lemma 4.3. Let $na \in \mathbb{N}a$ and let $\sum_k b_{k+1} q_k$ be the Ostrowski representation of n . Then

$$f(na) = \sum_k b_{k+1} \beta_k.$$

Proof. Let $m := \sum_k b_{k+1} p_k$. Then

$$na - m = \sum_k b_{k+1} (q_k a - p_k) = \sum_k b_{k+1} \beta_k \in I.$$

□

So in particular, $f(q_k a) = \beta_k$ for every $k \in \mathbb{N}$.

Corollary 4.4. The set $\{q_k a : k \text{ odd}\}$ is definable in \mathcal{R}_a .

Proof. Since $\beta_k < 0$ iff k is odd, we have by Lemma 4.3 that $f(q_k a) < 0$ iff k is odd. Hence the above set is equal to $\{na \in V : f(na) < 0\}$. □

Definition 4.5. Define $g : V \rightarrow \mathbb{R}^2$ by

$$q_l a \mapsto \begin{cases} (-\beta_l + \beta_{l+1}), -\beta_{l+1}), & \text{if } l \text{ is even,} \\ (-\beta_l, -(\beta_l + \beta_{l+1})), & \text{otherwise.} \end{cases}$$

By Lemma 4.3, Corollary 4.4 and $sv(q_l a) = q_{l+1} a$, the function g is definable. For ease of notation, we will write $g_l = (g_{l,1}, g_{l,2})$ for $g(q_l a)$.

Lemma 4.6. Let $n \in \mathbb{N}$ and $c \in \mathbb{R}$ be such that $\sum_{k=0}^{\infty} b_{k+1} \beta_k$ is the Ostrowski representation of c . If $-\beta_n < c < -(\beta_n + \beta_{n+1})$ or $-(\beta_n + \beta_{n+1}) < c < -\beta_n$, then $b_{k+1} = 0$ for all $k \leq n$.

Proof. Suppose $-\beta_n < c < -(\beta_n + \beta_{n+1})$. Since $\beta_{n+1} < 0$, n is even. By Fact 2.11 the Ostrowski representation of $-\beta_n$ is

$$a_{n+2} \beta_{n+1} + a_{n+4} \beta_{n+3} + a_{n+6} \beta_{n+5} + \dots$$

Since $-\beta_n < c$, we have $b_{k+1} = 0$ for all odd $k \leq n$ by Fact 2.13. By Fact 2.11 the Ostrowski representation of $-(\beta_n + \beta_{n+1})$ is

$$(a_{n+2} - 1) \beta_{n+1} + a_{n+3} \beta_{n+2} + a_{n+5} \beta_{n+4} + a_{n+7} \beta_{n+6} + \dots$$

Since $c < -(\beta_n + \beta_{n+1})$, we get $b_{k+1} = 0$ for all even $k \leq n$ by Fact 2.13. Hence $b_{k+1} = 0$ for all $k \leq n$. The case that $-(\beta_n + \beta_{n+1}) < c < -\beta_n$ can be handled similarly. \square

Lemma 4.7. Let $l, n \in \mathbb{N}$ and $c \in I$ such that $n < q_{l+1}$ and

$$f(na) + g_{l,1} \leq c < f(na) + g_{l,2}.$$

and let $\sum_{k=0}^{\infty} b_{k+1} \beta_k$ be the Ostrowski representation of c . Then $\sum_{k=0}^l b_{k+1} q_k$ is the Ostrowski representation of n .

Proof. Let l be even and $n \geq q_l$. Then by definition of the function g , we have

$$-(\beta_l + \beta_{l+1}) < c - f(na) < -\beta_{l+1}.$$

Hence by Lemma 4.6 the Ostrowski representation of $c - f(na)$ is $\sum_{k=l+1}^{\infty} c_{k+1} \beta_k$, for some $c_{k+1} \in \{0, \dots, a_{k+1}\}$. Now let $\sum_{k=0}^l c_{k+1} q_k$ be the Ostrowski representation of n . By Lemma 4.3

$$\sum_{k=0}^{\infty} b_{k+1} \beta_k = c = c - f(na) + f(na) = \sum_{k=l+1}^{\infty} c_{k+1} \beta_k + \sum_{k=0}^l c_{k+1} \beta_k = \sum_{k=0}^{\infty} c_{k+1} \beta_k.$$

Hence by Lemma 4.3 and the uniqueness of the Ostrowski representation, we have that $\sum_{k=0}^l b_{k+1} q_k = \sum_{k=0}^l c_{k+1} q_k$ is the Ostrowski representation of n . The case that l is odd can be shown similarly. \square

Lemma 4.8. Let $l \in \mathbb{N}$. For every $c \in I$ there is a unique $n \in \mathbb{N}_{< q_{l+1}}$ such that

$$(4.1) \quad f(na) + g_{l,1}(na) \leq c < f(na) + g_{l,2}(na).$$

Proof. Let $\sum_{k=0}^{\infty} b_{k+1} \beta_k$ be the Ostrowski representation of c . We first show the existence of such a natural number n . Set $n := \sum_{k=0}^l b_{k+1} q_k$. We will now show that n satisfies (4.1). Suppose that l is even. Since l is even, $\beta_{l+2k} > 0$ for each $k \in \mathbb{N}$. Then by Fact 2.11

$$f(na) - \beta_{l+1} = \sum_{k=0}^l b_{k+1} \beta_k + a_{l+3} \beta_{l+2} + a_{l+5} \beta_{l+4} + a_{l+7} \beta_{l+6} + \dots > c.$$

Suppose that $n \geq q_l$. Then we have $b_{l+1} > 0$. Hence by Fact 2.11

$$f(na) - (\beta_l + \beta_{l+1}) = \sum_{k=0}^l b_{k+1}\beta_k + (a_{l+2} - 1)\beta_{l+1} + a_{l+4}\beta_{l+3} + a_{l+6}\beta_{l+5} + \cdots \leq c.$$

Note that the inequality on the right follows immediately from $b_{l+1} > 0$. Now consider that $n < q_l$. Then by Fact 2.11

$$f(na) - \beta_l = \sum_{k=0}^l b_{k+1}\beta_k + a_{l+2}\beta_{l+1} + a_{l+4}\beta_{l+3} + a_{l+6}\beta_{l+5} + \cdots \leq c.$$

Hence (4.1) holds, if l is even. The case that l is odd can be treated similarly. The uniqueness of n follows directly from Lemma 4.7 and the uniqueness of Ostrowski representations. \square

Definition 4.9. Let $h : V \times I \rightarrow \mathbb{N}a$ map a pair $(q_l a, c)$ to the unique $na \in \mathbb{N}a_{< q_{l+1}a}$ given by Lemma 4.8.

Definition 4.10. We define

$$\begin{aligned} E_0 &:= \{(q_l a, c) \in V \times I : h(q_l a, c) < q_l a\}, \\ E_1 &:= \{(q_l a, c) \in V \times I : q_l a \leq h(q_l a, c) < \min\{q_{l+1}a, 2q_l a\}\}. \end{aligned}$$

Lemma 4.11. Let $i \in \{0, 1\}$, $l \in \mathbb{N}$, $c \in I$ and let $\sum_{k=0}^{\infty} b_{k+1}\beta_k$ be Ostrowski representation of c . Then $b_{l+1} = i$ iff $(q_l a, c) \in E_i$.

Proof. Let $q_l a \in V$. Then by Lemma 4.7, $\sum_{k=0}^l b_{k+1}q_k$ is the Ostrowski representation of $h(q_l a, c)/a$. Since the Ostrowski representation of a natural number is obtained by a greedy algorithm, we have $q_l \leq h(q_l a, c)/a < \min\{q_{l+1}, 2q_l\}$ iff $b_l = 1$, and $h(q_l a, c)/a < q_l$ iff $b_l = 0$. The statement of the Lemma follows immediately. \square

Definition 4.12. Define J to be set of $c \in I$ such that $(q_l a, c) \in E_0 \cup E_1$ for all $q_l a \in V$. Let $d \in J$ be the unique element in J such that $(a, d) \in E_1$ and

$$\forall q_l a \in V \ (q_l a, d) \in E_1 \text{ iff } (q_{l+1}a, d) \notin E_1.$$

Let $W := \{q_l \in V : (q_l a, d) \in E_1\}$.

It is easy to check that $q_l a \in W$ iff l is odd.

Definition 4.13. Define J' to be the set of all $c \in J$ such that $(q_l a, c) \in E_0$ whenever $q_l a \notin W$. Define $h_1 : W \rightarrow \mathbb{N}$ to be the function that maps $q_l a$ to $\frac{l-1}{2}$. Let $h_2 : J' \rightarrow \mathcal{P}(\mathbb{N})$ be the function that maps $c \in J'$ to $\{\frac{l-1}{2} : (q_l a, c) \in E_1\}$.

Theorem 4.14. The map $h = (h_1, h_2) : (W, J', s_W, E_1) \rightarrow (\mathbb{N}, \mathcal{P}(\mathbb{N}), s_{\mathbb{N}}, \in)$ is an isomorphism.

Proof. It follows immediately from the remark after the definition of W that $h_1 : (W, s_W) \rightarrow (\mathbb{N}, s_{\mathbb{N}})$ is an isomorphism. By definition of W , we have that $c \in J$ is in J' if and only if $(q_l a, c) \in E_0$ for every even $l \in \mathbb{N}$. Given a subset $X \subseteq \mathbb{N}$, one can easily find a unique $c \in I$ such that

$$c = \sum_{k \in X} \beta_{2k+1}.$$

We directly get that $c \in J'$ and for every $k \in \mathbb{N}$, we have $k \in X$ iff $(q_{2k+1}a, c) \in E_1$. Hence $h_2(c) = X$ and c is the unique element in J' with this property. From the construction it follows directly that $(q_la, c) \in E_1$ iff $h_1(q_la) \in h_2(c)$, for every $q_la \in W$ and $c \in J'$. \square

Hence \mathcal{R}_a defines an isomorphic copy of \mathcal{B} . This proves Theorem D.

5. DEFINING MULTIPLICATION IN \mathcal{R}_φ

Let $\varphi := \frac{1+\sqrt{5}}{2}$ be the golden ration. In this section it will be shown that multiplication by φ is definable in \mathcal{R}_φ . Since the continued fraction expansion of φ is $[1; 1, \dots]$, we get by Fact 2.6 that q_k is the k -th Fibonacci number, while p_k is $k+1$ -Fibonacci number. So in particular, $q_{k+1} = p_k$ and $\beta_k = q_k\varphi - q_{k+1}$. Moreover, because of the special form of the continued fraction expansion of φ , we get that $\zeta_k = \varphi$ for every $k \in \mathbb{N}$. Hence $\beta_{k+1} = -\frac{\beta_k}{\varphi}$ by Fact 2.7. Loosely speaking, this will allow us to realize multiplication by φ as a shift operation on the Ostrowski representations.

We will use the notation from the previous section. In particular, f, E_0 and E_1 are as defined before.

Definition 5.1. Let $L : \mathbb{N}\varphi \rightarrow \mathbb{N}\varphi$ map $n\varphi \in \mathbb{N}\varphi$ to the unique element $m\varphi \in \mathbb{N}\varphi$ such that $(q_k\varphi, f(m\varphi)) \in E_1$ iff $(q_{k+1}\varphi, f(n\varphi)) \in E_1$ for every $k \geq 1$. Let $T_1 : \mathbb{N}\varphi \rightarrow \mathbb{R}$ map $n\varphi$ to

$$\begin{cases} L(n\varphi) - f(L(n\varphi)) + 1, & \text{if } (\varphi, f(n\varphi)) \in E_1; \\ L(n\varphi) - f(L(n\varphi)), & \text{otherwise.} \end{cases}$$

Let $T_2 : \mathbb{N}\varphi \rightarrow \mathbb{R}$ map $n\varphi$ to

$$\begin{cases} f(L(n\varphi)) + \varphi - 1, & \text{if } (\varphi, f(n\varphi)) \in E_1; \\ f(L(n\varphi)), & \text{otherwise.} \end{cases}$$

Lemma 5.2. Let $n \in \mathbb{N}$ and let $\sum_k b_{k+1}q_k$ be the Ostrowski representation of n . Then the Ostrowski representation of $\varphi^{-1}L(n\varphi)$ is

$$\sum_k b_{k+2}q_k.$$

Proof. Set $m := \varphi^{-1}L(n\varphi)$ and let $\sum_k c_{k+1}q_k$ be the Ostrowski representation on m . It is left to show that $c_{k+1} = b_{k+2}$ for every $k \in \mathbb{N}$. By Lemma 4.3 the Ostrowski representation of $f(m\varphi)$ is $\sum_k c_{k+1}\beta_k$ and the Ostrowski representation of $f(n\varphi)$ is $\sum_k b_{k+1}\beta_k$. By Lemma 4.11, $c_{k+1} = 1$ iff $(q_k\varphi, f(m\varphi)) \in E_1$. By definition of the map L , this occurs iff $(q_{k+1}\varphi, f(n\varphi)) \in E_1$. Again, by Lemma 4.11 this happens if and only if $b_{k+2} = 1$. Hence $c_{k+1} = b_{k+2}$. \square

Lemma 5.3. Let $n \in \mathbb{N}$. Then $T_1(n\varphi) = n$.

Proof. Let $\sum_k b_{k+1}q_k$ be the Ostrowski representation of n . By Lemma 5.2 and $\beta_k = q_k\varphi - q_{k+1}$, we get that

$$\begin{aligned} n &= \sum_k b_{k+1}q_k \\ &= \left(\sum_{k>1} b_{k+1}q_{k-1}\right)\varphi + \left(\sum_{k>1} b_{k+1}(q_k - \varphi q_{k-1})\right) + b_2q_1 \\ &= L(n\varphi) - f(L(n\varphi)) + b_2 \\ &= T_1(n\varphi). \end{aligned}$$

□

Lemma 5.4. Let $n \in \mathbb{N}$. Then $\varphi f(n\varphi) = -f(T_2(n\varphi))$.

Proof. Since $\zeta_k = \varphi$ and $\beta_k = q_k\varphi - q_{k+1}$, we have by Fact 2.7 that

$$(5.1) \quad \frac{q_{k+1}\varphi - q_{k+2}}{q_k\varphi - q_{k+1}} = -\frac{1}{\varphi}.$$

Let $\sum_k b_{k+1}q_k$ be the Ostrowski representation of n . Note that

$$\varphi\beta_1 = \varphi(\varphi - 2) = \varphi^2 - 2\varphi = 1 - \varphi.$$

Hence by (5.1)

$$\begin{aligned} \varphi f(n\varphi) &= \sum_k b_{k+1}\varphi(q_k\varphi - q_{k+1}) = -\left(\sum_{k>0} b_{k+1}(q_{k-1}\varphi - q_k)\right) + \varphi b_2\beta_1 \\ &= -f(L(n\varphi) - b_2(\varphi - 1)) = -T_2(n\varphi). \end{aligned}$$

□

Theorem 5.5. The function $\lambda_\varphi : \mathbb{R} \rightarrow \mathbb{R}$ that maps $x \mapsto \varphi x$, is definable in \mathcal{R}_φ .

Proof. It is enough to define λ_φ on $\mathbb{R}_{\geq 0}$. For $m \in \mathbb{N}$ and $n\varphi \in \mathbb{N}\varphi$, define a map $P : \mathbb{N} \times \mathbb{N}\varphi \rightarrow \mathbb{R}$ by

$$P(m, n\varphi) := T_1^{-1}(m) - f(T_2(n\varphi)).$$

This is well-defined, since T_1 is injective by Lemma 5.3, and moreover definable in \mathcal{R}_φ . By Lemma 5.3 and Lemma 5.4, we have

$$P(m, f(n\varphi)) = T_1^{-1}(m) - f(T_2(n\varphi)) = \varphi m + \varphi f(n\varphi) = \varphi \cdot (m + f(n\varphi)).$$

Hence if there are $m, m' \in \mathbb{N}$ and $n\varphi, n'\varphi \in \mathbb{N}\varphi$ with $m + f(n\varphi) = m' + f(n'\varphi)$, we get $P(m, f(n\varphi)) = P(m', f(n'\varphi))$. Let $Q : \mathbb{N} + f(\mathbb{N}\varphi) \rightarrow \mathbb{R}$ map $m + f(n\varphi)$ to $P(m, f(n\varphi))$. By the above, Q is well-defined, definable in \mathcal{R}_φ and $Q(x) = \varphi x$ for all $x \in \mathbb{N} + f(\mathbb{N}\varphi)$. Since $\mathbb{N} + f(\mathbb{N}\varphi)$ is dense in $[1 - \varphi, \infty)$ and multiplication by φ is continuous, the graph of λ_φ on $[1 - \varphi, \infty)$ is the topological closure of the graph of Q in \mathbb{R}^2 . Hence λ_φ is definable in \mathcal{R}_φ . □

Theorem B now follows immediately from Theorem A and Theorem 5.5.

6. OPTIMALITY AND OPEN QUESTIONS

1. Let \mathcal{L} be the language of \mathcal{R}_a for some $a \in \mathbb{R}$. For $a \in \mathbb{R} \setminus \mathbb{Q}$, we have seen that the structure \mathcal{R}_a defines the set $\{q_k a : k > 0\}$, which we denoted by V . Since $s_V^{k-1}(a) = q_k a$, it is easy to see that for every $k, l \in \mathbb{N}$ there is an \mathcal{L} -sentence $\psi_{k,l}$ such that for all $a \in \mathbb{R} \setminus \mathbb{Q}$

$$\mathcal{R}_a \models \psi_{k,l} \text{ iff } q_{k+l} = lq_k + q_{k-1}.$$

It follows immediately from Fact 2.6 that if $a = [a_0; a_1, \dots]$ and the function that takes k to a_k is non-computable, then the theory of \mathcal{R}_a is undecidable.

2. For a quadratic, quantifier elimination results for \mathcal{B} like [4, Theorem 1] transfer quite directly to \mathcal{R}_a because of Theorem C and Theorem D. Any attempt of proving substantially different quantifier elimination results for \mathcal{R}_a are likely to fail due to Theorem D.

3. Let $a \in \mathbb{R} \setminus \mathbb{Q}$. By Lemma 4.3, the function $f : \mathbb{N}a \rightarrow \mathbb{R}$ that takes $na \in \mathbb{N}a$ to $\sum_k b_k \beta_k$, where $\sum_k b_k q_k$ is the Ostrowski representation of n , is definable in \mathcal{R}_a . This function maps a closed and discrete set onto a dense subset of the interval $[1-a, 2-a)$. Hence together with Theorem A of the current paper, it follows that for a quadratic the structure \mathcal{R}_a satisfies condition (i) of [8, Theorem A], but not its conclusion. Hence condition (ii) can not be dropped from [8, Theorem A].

4. Except for Theorem D not much is known about the structure \mathcal{R}_a when a is not quadratic. For example it is not known whether there is an a such that \mathcal{R}_a defines multiplication on \mathbb{R} . Even in the case of Euler's number e we do not know whether the theory of \mathcal{R}_e is decidable or not. Because the continued fraction expansion of e is not periodic, it is unlikely that \mathcal{R}_e can be defined in \mathcal{B} , surely not in the way presented here. On the one hand the continued fraction expansion of e is simple enough that other methods might be used to show decidability, but on the other hand the expansion S_e defines multiplication on \mathbb{R} by [8, Theorem B].

5. It is an open question whether Theorem B holds for all quadratic numbers. However, it seems possible that a closer examination of how multiplication by a interacts with the Ostrowski representations based on a , might at least yield decidability of the theory of \mathcal{S}_a for every quadratic a .

6. By [8, Theorem C], $(\mathbb{R}, <, +, \mathbb{Z}, a\mathbb{Z}, b\mathbb{Z})$ defines multiplication on \mathbb{R} and hence every projective set whenever $1, a, b$ are linearly independent over \mathbb{Q} . However, we do not know whether there is a set definable in both $(\mathbb{R}, <, +, \mathbb{Z}, a\mathbb{Z})$ and $(\mathbb{R}, <, +, \mathbb{Z}, b\mathbb{Z})$ that is not definable in $(\mathbb{R}, <, +, \mathbb{Z})$. It would be interesting if a version of Cobham's theorem similar to the results in Boigelot, Brusten and Bruyère [2] holds in this setting.

7. Let $a \in \mathbb{R} \setminus \mathbb{Q}$. Note that an isomorphic copy of \mathcal{R}_a is definable in the expansion $(\mathbb{R}, <, +, \cdot, e^{\mathbb{Z}}, e^{\mathbb{Z}a})$ of the real field, but by [6, Theorem 1.3] the theory of the latter structure is undecidable, even if a is quadratic.

8. Several of the results in this paper can be reformulated to state that certain sets are recognizable by certain automata. For example, Lemma 3.15 shows that the graph of addition of real numbers given in Ostrowski representation based on a quadratic irrational is recognizable by a deterministic Müller automaton. While not crucial for the main results of this paper, it might still be interesting to give and study explicit constructions of these automata.

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